On the Approximability of Weighted Model Integration over DNF Structures [1]

Ralph Abboud, İsmail İlkan Ceylan, Radoslav Dimitrov







propositional formula

 ϕ $w: \mathcal{A}(\phi) \to \mathbb{R}$

propositional formula weight function

 ϕ

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 $WMC(\phi)$

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Applications

- Probabilistic Graphical Models
- Probabilistic Databases

propositional formula

weight function



Probabilistic Logic Programming

Probabilistic Knowledge Bases

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X set of *n* real variables V set of *m* Boolean variables

 $c_1 x_1 + \ldots + c_i x_i \bowtie c$

LRA Atom, $x_i \in X$, $\bowtie \in \{<, \leq, >, \geq, =, \neq\}$

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atoms(X, V)

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propositional formula over $X \cup V$

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$c_1 x_1 + \ldots + c_i x_i \bowtie c$	LRA Atom, $x_i \in X$, $\bowtie \in \{<, \le, >, \ge, =, \neq\}$
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$w: \mathbb{R}^n \times \mathbb{B}^m \to \mathbb{R}$	weight function over $X \cup V$

 $\sum_{\nu} \int_{x_{\phi}} w(x,\nu) \, dx$

 $WMI(\phi)$

where \boldsymbol{v} is a Boolean assignment over V, x_{ϕ} denotes valuations of X satisfying ϕ .

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For weight functions in WMI, it is common to factorize [2] *w* as a product of *m* Boolean literal weights and a density function over real variables, i.e.,:

$$w(x,v) = w_x(x) \prod_{i=1}^m w_b(p_i).$$

 $(x \lor y) \land (y \lor z)$

Conjunctive Normal Form (CNF)

 $(x \lor y) \land (y \lor z)$

 $(x \wedge y) \vee (y \wedge z)$

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Conjunctive Normal Form (CNF)

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Special cases WMI(CNF), WMI(DNF)

 $(x \lor y) \land (y \lor z)$

 $(x \land y) \lor (y \land z)$

Conjunctive Normal Form (CNF)

Disjunctive Normal Form (DNF)



Both WMI and WMC are **#P-hard** for exact solving. Hence, we study WMI within the context of approximate solving.

Approximation Hardness

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	WMC		WMI
CNF	NP-hard [3]		NP-hard
DNF	FPRAS [4]	➡	?

Result:

Show that WMI(DNF) admits an FPRAS for concave weight functions

Result builds on existing FPRAS algorithms for WMC(DNF) and volume computation for the union of convex bodies

Consider *k* convex bodies. We wish to compute the volume of their union.



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Run until Step 4 executes T times:

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Compute individual body volumes

Run until Step 4 executes T times:



Randomly sample body

sampling probabilities proportional to body volume

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Uniformly sample another body

If sampled point in this body, Success! Repeat Step 2, Else repeat Step 4.

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ApproxUnion: Volume of Union of Convex Bodies [5]

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Number of successes then yields an unbiased estimator for volume

ApproxUnion is an FPRAS



Volume Computation: FPRAS with error ϵ_V , confidence δ_V



Point Sampling: FPRAS with error ϵ_S , confidence δ_S



Membership Check: FPRAS with error ϵ_P , confidence δ_P

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ApproxUnion is an FPRAS with error ε and confidence δ , using $T = O(k\varepsilon^{-2})$, for $\varepsilon_V, \varepsilon_S \leq \frac{\varepsilon^2}{47k}, \varepsilon_P \leq \frac{\varepsilon^2}{47k^2}, \delta_V \leq \frac{\delta}{4k}, \delta_S + \delta_P \leq \frac{\delta}{2276 \ln(\frac{8}{\delta})\frac{k}{\varepsilon^2}}$.

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- ApproxWMI:
 - 1) Computes the weight of every DNF clause: weight function integral
 - over LRA-induced convex polytope, multiplied by Boolean probabilities.
 - 2) Samples points from a clause as per the weight function.
 - 3) Checks membership of point to another uniformly random clause.

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Concave Weight Function



2D Convex Polytope

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- The volume of this polytope is then computed using standard tools [6].



2D Convex Polytope

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- The volume of this polytope is then computed using standard tools [6].
- This volume is multiplied by Boolean probabilities to return clause weight.



- The *n+1*-dimensional polytope is sampled uniformly, as a proxy for sampling the original polytope according to the weight function.



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- Membership to LRA polytope is checked by validating point against every LRA atom in a clause.





ClauseWeight: FPRAS with error ϵ_V , confidence δ_V



Sampling: FPRAS with error ϵ_S , confidence δ_S



Evaluate: FPRAS with error ϵ_P , confidence δ_P



ClauseWeight: FPRAS with error ϵ_V , confidence δ_V



3 Evaluate: FPRAS with error ϵ_P , confidence δ_P

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- Number of clauses is roughly the number of variables divided by clause width.
- Real weight function is a concave polynomial with degree up to 5.

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- Target error *ε*: Set to 0.15, 0.25, and 0.35.

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Three algorithm configurations:

- Target error *ε*: Set to 0.15, 0.25, and 0.35.
- Target confidence δ : Set to 0.05, 0.15 and 0.25.

Results


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- ApproxWMI is scalable, and can efficiently solve instances with up to 1K variables, which is out of reach for existing WMI solvers.
- ApproxWMI can be extended to more general factorizations enabling real-Boolean dependency.
- ApproxWMI is a useful tool for efficient probabilistic inference in hybrid domains.

Thank You!



Selected References

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